

A WEIGHTED TURÁN SIEVE FOR TWIN PRIMES VIA THE KRAFFT GEOMETRY

FERNANDO PORTELA 

Ronin Institute for Independent Scholarship 2.0, USA

ABSTRACT. The Twin Prime Conjecture is notoriously obstructed by the parity barrier, which prevents classical multiplicative sieves from isolating prime pairs. In this paper, we introduce an additive Turán sieve framework evaluated over bounded, symmetric intervals. Utilizing a centered modular alignment (which we term the Krafft alignment), we construct an additive sieve formulation where the existence of twin primes corresponds to the vanishing of a local penalty function. By evaluating the character sums associated with this sieve, we isolate a destructive interference term at the evaluation point $h = 3q/p$. We reduce the twin prime problem to a variational optimization problem over a finite-dimensional parameter space: the existence of twin primes in a prescribed interval is guaranteed, provided the minimum sieve weight quotient satisfies $\mu_{\min}(n) < 1$. We further show that independent, one-dimensional sieve weights are inherently insufficient to cross this barrier, structurally requiring the use of multidimensional correlations to avoid the quotient $\mu \geq 1$. This structural deficit constitutes a spectral manifestation of the Selberg parity barrier, demonstrating that any successful sieve must intrinsically rely on higher-order correlations and multidimensional exponential sums. All discrete definitions and structural lemmas in this paper have been formally verified in Lean 4, including the conditional implication that if $\mu_{\min}(n) < 1$ holds for infinitely many n , then there are infinitely many twin primes.

1. INTRODUCTION

The twin prime conjecture is notoriously obstructed by the parity barrier, which prevents classical sieve methods from isolating prime pairs. In this paper, we develop a new additive Turán sieve framework [3, 20] designed to bypass this combinatorial obstruction. While classical methods rely on inclusion-exclusion principles to track surviving candidates, the Turán sieve instead bounds the second moment (the variance) of an additive penalty function. In the classical Turán variance method, this penalty function typically counts the number of distinct prime factors of an integer. Here, we adapt this methodology by evaluating the variance of a weighted hit-count of forbidden congruences over a highly structured, symmetric interval. This is achieved through a centered modular alignment—which we term the Krafft alignment [12, 16]. We guarantee the existence of twin primes by establishing a positive variance gap—a condition that minimizes the sieve quotient $\mu_{\min}(n)$ —ensuring the main term dominates the resonant error terms. By analyzing the

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spectral properties of this penalty function, we reduce the twin prime conjecture to a finite-dimensional variational optimization problem.

1.1. The Parity Barrier and Sieve Limitations. To contextualize the motivation for this approach, we must consider the historical context of the parity barrier. Classical benchmarks for prime density were established by the Hardy-Littlewood conjecture [11]. Subsequent sieve methods, originating with Brun [2] and refined by Selberg [17] and others [10, 14], have provided the main avenues of attack on prime distribution. While these frameworks have yielded extraordinary structural results—most notably the bounded gaps between primes [4, 5, 8, 13, 21]—they inevitably encounter Selberg’s parity barrier. This phenomenon demonstrates that classical sieves—whose weights depend fundamentally on the Möbius function or similarly bounded multiplicative functions—cannot distinguish between integers possessing an even number of prime factors and those possessing an odd number. Because the Möbius function fluctuates in sign based on this parity, these sieves structurally fail to isolate primes. This preclusion prevents a direct proof of the twin prime conjecture using classical multiplicative methods alone.

The term “parity barrier” functions much like “dark matter” in astrophysics: it is a name given to the *symptoms* of an inherent obstruction in multiplicative number theory, rather than a physical law. From an analytic and multiplicative perspective, as elucidated by Friedlander and Iwaniec [7], this barrier is inextricably tied to the analytic behavior of generating functions (such as the multiple-order poles of the Dirichlet series for higher-order von Mangoldt functions Λ_k) and the possible existence of exceptional zeros.

In this paper, we provide a *spectral perspective* on this obstruction. When our additive sieve is restricted to independent, one-dimensional weights, the parity barrier manifests spectrally as a destructive interference at specific rational frequencies (e.g., $h = 3q/p$). Because isolated character sum evaluations cannot outpace the logarithmic divergence of Mertens’ sums, the sieve quotient is forced to $\mu \geq 1$. This analytic deficit reveals that resolving the conjecture requires exploiting the full multidimensional space of higher-order correlations, a domain where this obstruction can be systematically bypassed.

1.2. The Krafft Geometry and the Sieve Domain. To circumvent the parity limitations inherent in multiplicative sieves, we transition to an additive framework evaluated over a bounded, symmetric interval. The underlying geometric transformation we employ is rooted in an arithmetic equivalence first noted by the astronomer W. L. Krafft in 1798 [12]. Krafft, who had been Euler’s assistant at the Imperial Academie of Sciences in St Petersburg, observed that finding twin primes is equivalent to finding natural numbers m that cannot be expressed in the Diophantine form $m = 6ab \pm a \pm b$. While this specific form remains relatively obscure (yet noted by Sierpiński [18] and documented as the OEIS sequence A002822 [15]), it can be algebraically reframed as a parametric linear construct: $m = (6b \pm 1)a \pm b$.

By fixing the parameter $b \in \mathbb{N}$, this condition generates an infinite set of arithmetic progressions, which can be visualized as a grid (Figure 1) where rows correspond to the moduli $(6b \pm 1)$, justifying our first definition:

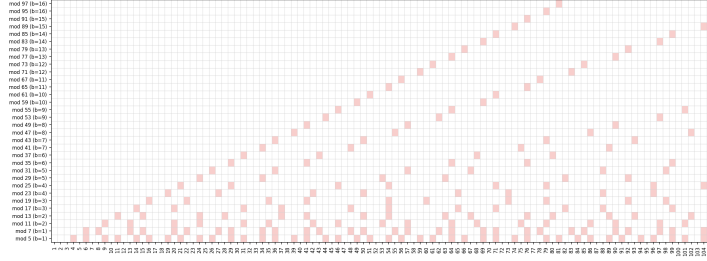


FIGURE 1. Visualization of the Krafft arithmetic progressions.

Definition 1.1 (The Krafft Residues $[K]$). For each prime $p_i \in \mathcal{P}_n$, we establish the target Krafft alignment residue as:

$$r_i^K = \left\lfloor \frac{p_i + 1}{6} \right\rfloor$$

A simple high-school level observation reveals that rows corresponding to composite moduli are completely redundant (Figure 2), meaning we only need to sift against prime moduli $p \geq 5$.

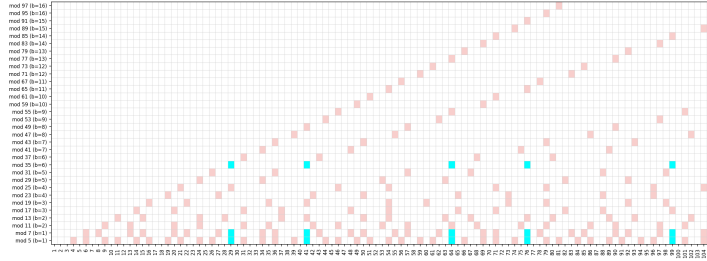


FIGURE 2. Redundancy of composite moduli in the Krafft grid. Depicted: (mod 35) "collides" with both (mod 5) and (mod 7)

Furthermore, the grid exhibits a profound symmetry (Figure 3) that naturally bounds the search space, giving rise to Evaluation Intervals within which twin prime candidate indices must reside (Figure 4).

To translate this visual symmetry into a rigorous analytic framework, we formally define the prime window and the evaluation domain.

Definition 1.2 (The Prime Window and Primorial $[P]$ & $[P]$). Let \mathcal{P}_n denote the set of primes evaluated by the sieve, bounded as:

$$\mathcal{P}_n = \{p \text{ prime} \mid 5 \leq p < 6n + 2\}$$

We define $w(n) = |\mathcal{P}_n|$ as the total count of these prime factors, and $q = \prod_{p \in \mathcal{P}_n} p$ as their associated primorial.

Definition 1.3 (The Evaluation Interval $[E]$). We constrain our search for twin prime indices to the highly structured symmetric interval \mathcal{A}_n :

$$\mathcal{A}_n = [6n^2 - 2n, 6n^2 + 10n + 3]$$

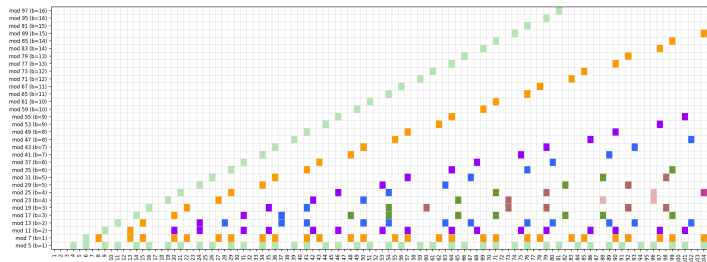


FIGURE 3. Shadowing in the Krafft grid. The first forbidden residues of all higher moduli form a pseudo-diagonal that is perfectly “eclipsed” by the constraints of the lowest prime (mod 5, light green). Similarly, the second forbidden residues are entirely hidden behind the constraints of the next prime (mod 7, orange). This sequential shadowing renders infinite families of higher-moduli constraints strictly redundant.

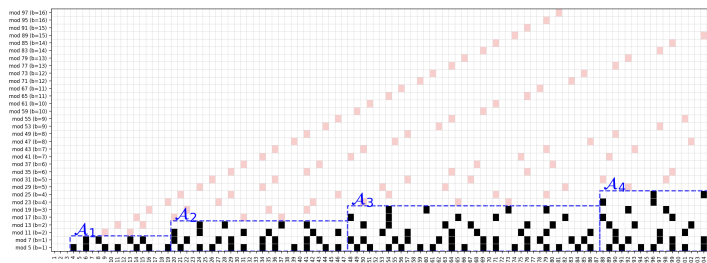


FIGURE 4. Evaluation intervals. Removing the redundant pseudo-diagonal alignments leaves clearly delineated, linearly growing in size domains (formalized in Definition 1.3).

Remark 1.4. This geometric setup readily yields a heuristic estimate for the density of twin prime indices within our evaluation interval. By mapping the twin prime condition to the avoidance of two specific residue classes modulo each prime p , the proportion of surviving candidates is governed by the product $\prod_{p \in \mathcal{P}_n} (1 - 2/p)$. By Mertens’ theorems, this product is asymptotic to $O(1/\log^2 x)$. Assuming the surviving candidates are well-distributed, an interval of size x is thus heuristically expected to contain $O(x/\log^2 x)$ twin prime pair indices, perfectly mirroring the classical Hardy-Littlewood prediction.

1.3. Outline and Formalization. The paper is organized as follows. Section 2 completes the additive sieve equivalence. Section 3 shifts the problem into the frequency domain, evaluating the necessary discrete Fourier transforms and identifying the structural destructive interference. Section 4 proves the spectral manifestation of the parity barrier, showing that one-dimensional weights are insufficient. Section 5 establishes the continuous variational framework and proves that the minimization of our sieve quotient $\mu_{\min}(n)$ guarantees the existence of twin primes. Finally, Section

6 explores the full character convolution and presents numerical optimizations that empirically validate the structural necessity of multidimensional correlations.

Throughout the text, [↗](#) symbols are used as direct links to formalized code in the GitHub repository github.com/ElNando888/KrafttSieve/tree/v1.8, which contains the machine-checked verifications for all definitions, statements, and proofs in this paper.¹ We use Lean [6] version 4.31.0-rc1 together with the Lean mathematical library `Mathlib` [19] revision `d568c8c`.

2. THE ADDITIVE SIEVE EQUIVALENCE

The specific interval width of \mathcal{A}_n is chosen so that the alignment residue r_i^K effectively translates the standard divisibility conditions of twin prime candidates into a single modular congruence evaluated at the index x . The following lemma makes this equivalence precise.

Lemma 2.1 (Modular Alignment Equivalence [[↗](#)]). *For any prime factor $p \geq 5$ and any integer x , let $r = \lfloor (p+1)/6 \rfloor$. Then $x \equiv \pm r \pmod{p}$ if and only if $p \mid (6x-1)$ or $p \mid (6x+1)$.*

Proof. Since $p \geq 5$ is prime, we have either $p \equiv 1 \pmod{6}$ or $p \equiv 5 \pmod{6}$. If $p \equiv 1 \pmod{6}$, then $r = (p-1)/6$, and therefore $6r = p-1 \equiv -1 \pmod{p}$. If $p \equiv 5 \pmod{6}$, then $r = (p+1)/6$, yielding $6r = p+1 \equiv 1 \pmod{p}$. In both cases, $6r \equiv \pm 1 \pmod{p}$. Assuming $x \equiv \pm r \pmod{p}$, multiplying both sides by 6 gives $6x \equiv \pm 6r \equiv \pm 1 \pmod{p}$. This implies $6x \mp 1 \equiv 0 \pmod{p}$, meaning p divides either $6x-1$ or $6x+1$. The steps are reversible, establishing the equivalence. \square

To track the sieve's progress additively rather than through a multiplicative inclusion-exclusion product, we establish an indicator function for the forbidden residue classes and a global penalty (or counting) function that accumulates across all primes in our window.

Definition 2.2 (Local Indicator and Global Penalty Function [[↗](#) & [↗](#)]). *We define the local indicator function $g_i(x) \in \{0, 1\}$ to detect whether an integer x falls into the forbidden residue classes modulo p_i :*

$$g_i(x) = \begin{cases} 1, & \text{if } x \equiv \pm r_i^K \pmod{p_i}; \\ 0, & \text{otherwise.} \end{cases}$$

The global penalty function $c(x)$ is defined additively as the sum of all local indicators across the prime window:

$$c(x) = \sum_{i=1}^{w(n)} g_i(x)$$

Unlike classical inclusion-exclusion setups which track the surviving elements via alternating signs, $c(x)$ is non-negative and simply counts the number of forbidden congruences satisfied by x . An integer with zero penalty ($c(x) = 0$) has successfully avoided all prime factors in \mathcal{P}_n . To ensure that avoiding these specific prime factors is sufficient to guarantee true primality, we rely on the strict algebraic boundaries of the interval \mathcal{A}_n .

¹The only axioms used are `propext`, `Classical.choice`, and `Quot.sound`, which can be checked by the `#print axioms` command.

Lemma 2.3 (Interval Projection Bound [43]). *For any integer $x \in \mathcal{A}_n$, the value $6x + 1$ is strictly less than the square of the next possible prime candidate after the sieve window, namely $6x + 1 < (6n + 5)^2$.*

Proof. The maximum value in the interval \mathcal{A}_n is $x = 6n^2 + 10n + 3$. Evaluating $6x + 1$ at this maximum yields $6(6n^2 + 10n + 3) + 1 = 36n^2 + 60n + 19$. The next prime candidate strictly greater than the upper bound of \mathcal{P}_n (which is $6n + 1$) is at least $6n + 5$. Comparing our maximal projection to the square of this candidate gives $(6n + 5)^2 = 36n^2 + 60n + 25$. Since $36n^2 + 60n + 19 < 36n^2 + 60n + 25$, the strict inequality $6x + 1 < (6n + 5)^2$ holds for all $x \in \mathcal{A}_n$. \square

This bound is key to the reduction, as it ensures that any survivors within the interval are forced to be prime.

Theorem 2.4 (Additive Sieve Formulation [43]). *For any integer $n \geq 1$ and any $x \in \mathcal{A}_n$, the global penalty function evaluates to exactly zero, $c(x) = 0$, if and only if both $6x - 1$ and $6x + 1$ are prime numbers.*

Proof. By the definition of the global penalty function, $c(x) = 0$ if and only if $g_i(x) = 0$ for all $1 \leq i \leq w(n)$. This means $x \not\equiv \pm r_i^K \pmod{p_i}$ for all primes $p_i \in \mathcal{P}_n$. By Lemma 2.1, this guarantees that neither $6x - 1$ nor $6x + 1$ has any prime divisors in \mathcal{P}_n . Integers of the form $6x \pm 1$ are intrinsically coprime to 2 and 3. Therefore, $6x - 1$ and $6x + 1$ have no prime factors strictly less than $6n + 5$. Because Lemma 2.3 structurally guarantees that the maximum possible value of $6x + 1$ across the entire interval \mathcal{A}_n is strictly less than $(6n + 5)^2$, any composite number of the form $6x \pm 1$ must have a prime factor strictly less than $6n + 5$. Since all such factors have been eliminated, it follows that both $6x - 1$ and $6x + 1$ must be prime. \square

3. CHARACTER SUM EXPANSION AND THE MAIN SIEVE IDENTITY

Historically, the pure Turán sieve, while elegantly replacing inclusion-exclusion mechanics with a second-moment variance bound, is relatively weak. Used alone, an unweighted variance bound struggles to strictly isolate prime numbers, which is why the method often appears only as a historical side-chapter in modern analytic number theory textbooks. The landscape of sieve theory was transformed when Selberg [17] introduced the groundbreaking concept of applying arbitrary, optimizable weights. This innovation shifted the sieve problem from rigid combinatorial accounting to a flexible analytic optimization problem.

Following this philosophy, we analyze the distribution of twin prime indices $x \in \mathcal{A}_n$ that avoid the local penalties, by constructing a weighted sifting condition in the spatial domain, which we will subsequently evaluate through its character sums over the cyclic group. We begin by introducing an arbitrary, real-valued weight function $W(x)$ strictly supported within our bounded evaluation interval \mathcal{A}_n .

Definition 3.1 (Sieve Weight Sums). *For an arbitrary real-valued weight function $W(x)$ supported strictly on \mathcal{A}_n , we define the total weight sum S_1 and the penalized weight sum S_2 :*

$$S_1(n, W) = \sum_{x \in \mathcal{A}_n} W(x)$$

$$S_2(n, W) = \sum_{x \in \mathcal{A}_n} W(x)c(x)$$

Lemma 3.2 (The Weighted Existence Principle [vri]). *If there exists a non-negative weight function $W(x)$ such that the global penalized sum is strictly less than the base weight sum, $S_2(n, W) < S_1(n, W)$, then there must exist at least one integer $x \in \mathcal{A}_n$ such that $W(x) > 0$ and $c(x) = 0$.*

Proof. The global penalty function $c(x)$ is a sum of non-negative indicator functions, forcing $c(x) \geq 0$. Consequently, any non-zero penalty must satisfy $c(x) \geq 1$. If $c(x) \geq 1$ for all coordinates where $W(x) > 0$, then the sum $S_2(n, W) = \sum W(x)c(x)$ must be greater than or equal to $S_1(n, W) = \sum W(x)$. This contradicts the strict inequality $S_2 < S_1$. Therefore, there must exist at least one index x satisfying $c(x) = 0$ and carrying strictly positive weight. \square

Definition 3.3 (Sieve Sufficiency Condition). *We say the sieve sufficiency condition holds for a given n if there exists a weight function $W : \mathbb{Z}/q\mathbb{Z} \rightarrow \mathbb{R}$ that is non-negative everywhere, supported on \mathcal{A}_n , and satisfies the strict inequality*

$$S_2(n, W) < S_1(n, W)$$

Theorem 3.4 (Sieve Sufficiency Guarantee [vri]). *For any integer $n \geq 1$, if Definition 3.3 is satisfied, then there exists an integer $x \in \mathcal{A}_n$ such that $6x - 1$ and $6x + 1$ form a twin prime pair.*

Proof. If Definition 3.3 is satisfied, then Theorem 3.2 guarantees the presence of at least one index $x \in \mathcal{A}_n$ satisfying zero penalty ($c(x) = 0$). Under Theorem 2.4, the vanishing of the penalty implies that both $6x - 1$ and $6x + 1$ are prime. \square

In the spatial domain, analyzing the pseudo-random overlaps of the forbidden residue classes across all primes is analytically intractable. We therefore shift into the frequency domain, where the strict periodicities of the local indicator functions resolve into highly sparse, orthogonal spectra, allowing us to separate the principal density from the oscillatory interference. This is accomplished by expanding our functions over the characters of the cyclic group $\mathbb{Z}/q\mathbb{Z}$.

Definition 3.5 (Fourier transform over $\mathbb{Z}/q\mathbb{Z}$). *For any function $f(x)$, the frequency coefficients of its Fourier transform over $\mathbb{Z}/q\mathbb{Z}$ are given by:*

$$\hat{f}(h) = \frac{1}{q} \sum_{x=0}^{q-1} f(x) e^{-2\pi i h x / q}$$

By exploiting the bounds of \mathcal{A}_n , we can extend our spatial sums over the entire cyclic group without altering their values. This allows us to apply Plancherel's theorem to expand the penalized mass into a series of exponential sums.

Lemma 3.6 (Compact Support Equivalence [vri]). *Assume the weight function $W(x)$ is supported within the interval \mathcal{A}_n , meaning $W(x) = 0$ for all $x \notin \mathcal{A}_n$. Then the total weight $S_1(n, W)$ is proportional to the zero-frequency Fourier coefficient of the weight function:*

$$S_1(n, W) = q \hat{W}(0)$$

Proof. Since $W(x)$ is zero outside \mathcal{A}_n , the sum over the interval is equivalent to the sum over the entire cyclic group $\mathbb{Z}/q\mathbb{Z}$. Evaluating Definition 3.5 at $h = 0$ yields

$$\hat{W}(0) = \frac{1}{q} \sum_{x \in \mathbb{Z}/q\mathbb{Z}} W(x) e^0 = \frac{1}{q} S_1(n, W)$$

Multiplying by q completes the proof. \square

Lemma 3.7 (The Character Sum Penalty Expansion [44]). *For a weight function $W(x)$ supported on \mathcal{A}_n , the weighted penalty sum $S_2(n, W)$ decomposes into the Fourier domain as:*

$$S_2(n, W) = q \sum_{i=1}^{w(n)} \sum_{h=0}^{q-1} \hat{W}(h) \overline{\hat{g}_i(h)}$$

Proof. Substitute the additive definition $c(x) = \sum g_i(x)$ into $S_2(n, W)$. Because $W(x)$ has compact support on \mathcal{A}_n , the sum extends to $\mathbb{Z}/q\mathbb{Z}$ without altering the value. Applying the standard discrete Plancherel theorem to $W(x)$ and $c(x)$, and utilizing the linearity of the Fourier transform over the finite sum of local indicators, yields the expansion. \square

Because the local indicator function $g_i(x)$ depends only on the residue of x modulo the local prime p_i , projecting it onto the full cyclic group $\mathbb{Z}/q\mathbb{Z}$ yields a highly sparse spectrum.

Lemma 3.8 (Indicator function spectrum [44 & 45]). *The discrete Fourier transform (as defined in Definition 3.5) of the local indicator function, $\hat{g}_i(h)$, evaluates to zero for all frequencies except those that are precise integer multiples of q/p_i . Specifically, let $h \in \mathbb{Z}/q\mathbb{Z}$:*

- (1) *If $h = k(q/p_i)$ for some integer $0 \leq k < p_i$, then*

$$\hat{g}_i\left(k \frac{q}{p_i}\right) = \frac{2}{p_i} \cos\left(\frac{2\pi k r_i^K}{p_i}\right)$$

- (2) *If h is not a multiple of q/p_i , then $\hat{g}_i(h) = 0$.*

Proof. By writing $x = y + mp_i$ with $0 \leq y < p_i$ and $0 \leq m < q/p_i$, we can factor the Fourier transform as

$$\begin{aligned} \hat{g}_i(h) &= \frac{1}{q} \sum_{x=0}^{q-1} g_i(x) e^{-2\pi i h x / q} \\ &= \frac{1}{q} \sum_{y=0}^{p_i-1} g_i(y) e^{-2\pi i h y / q} \sum_{m=0}^{q/p_i-1} e^{-2\pi i h m p_i / q} \end{aligned}$$

The inner sum over m vanishes identically unless h is a multiple of q/p_i . When $h = k(q/p_i)$, the sum over m evaluates to q/p_i , yielding

$$\hat{g}_i\left(k \frac{q}{p_i}\right) = \frac{1}{p_i} \sum_{y=0}^{p_i-1} g_i(y) e^{-2\pi i k y / p_i}.$$

Since $g_i(y) = 1$ only at the two residues $y \equiv \pm r_i^K \pmod{p_i}$, we obtain

$$\begin{aligned} \hat{g}_i\left(k \frac{q}{p_i}\right) &= \frac{1}{p_i} \left(e^{-2\pi i k r_i^K / p_i} + e^{2\pi i k r_i^K / p_i} \right) \\ &= \frac{2}{p_i} \cos\left(\frac{2\pi k r_i^K}{p_i}\right) \end{aligned}$$

\square

This sparsity allows us to explicitly separate the principal expected density term from the targeted oscillatory interference, yielding the main identity of the sieve.

When evaluating the exponential sum at $h = 3q/p_i$, the trigonometric phase aligns with the alignment residue, causing destructive interference.

Lemma 3.9 (Evaluation of the Character Sum at $h = 3q/p_i$ [vri]). *For any prime factor $p_i \geq 5$, substituting the alignment residue $r_i^K = \lfloor (p_i + 1)/6 \rfloor$ evaluated at $k = 3$ yields a negative term:*

$$\cos\left(\frac{2\pi \cdot 3 \cdot r_i^K}{p_i}\right) = -\cos\left(\frac{\pi}{p_i}\right)$$

Proof. Resolving the floor function $r_i^K = \lfloor (p_i + 1)/6 \rfloor$ gives $r_i^K = \frac{p_i \pm 1}{6}$, uniquely determined by the parity of $p_i \pmod{6}$. Evaluating the phase angle generates

$$\frac{6\pi}{p_i} \left(\frac{p_i \pm 1}{6}\right) = \pi \pm \frac{\pi}{p_i}$$

Applying the standard trigonometric identity $\cos(\pi \pm \theta) = -\cos(\theta)$ directly yields the result. \square

Theorem 3.10 (The Main Sieve Identity [vri]). *For any weight function $W(x)$ supported on \mathcal{A}_n , the penalized sum isolates into a principal expected density term and an oscillatory interference sum:*

$$S_2(n, W) = S_1(n, W) \sum_{i=1}^{w(n)} \frac{2}{p_i} + q \sum_{i=1}^{w(n)} \sum_{k=1}^{p_i-1} \hat{W}\left(k \frac{q}{p_i}\right) \frac{2}{p_i} \cos\left(\frac{2\pi k r_i^K}{p_i}\right) \quad (1)$$

Proof. Starting from Lemma 3.7, we partition the inner sum over frequencies h into the zero frequency $h = 0$ and the non-zero frequencies. Applying Lemma 3.8, all frequencies h that are not multiples of q/p_i evaluate to $\hat{g}_i(h) = 0$ and are discarded. The sum collapses entirely to the frequencies $h = k(q/p_i)$. Finally, applying Lemma 3.6 replaces $q\hat{W}(0)$ with the total weight $S_1(n, W)$, establishing the exact equation. \square

This destructive interference can suppress $S_2(n, W)$ below the total weight $S_1(n, W)$.

Lemma 3.11 (Isolation of the $h = 3q/p_i$ Term [vri]). *For a weight function $W(x)$ whose spectral support is restricted to the base harmonic $h = 0$ and the precise frequencies $h = 3q/p_i$, the sieve identity (1) simplifies exactly to:*

$$S_2(n, W) = S_1(n, W) \sum_{i=1}^{w(n)} \frac{2}{p_i} - q \sum_{i=1}^{w(n)} \hat{W}\left(3 \frac{q}{p_i}\right) \frac{2}{p_i} \cos\left(\frac{\pi}{p_i}\right)$$

Proof. By the spectral restriction of W , all terms in equation (1) where $k \neq 3$ vanish identically. For the surviving $k = 3$ terms, applying Lemma 3.9 inverts the sign of the interference sum, replacing the general cosine term with $-\cos(\pi/p_i)$. \square

4. THE PARITY BARRIER IN ONE DIMENSION

Definition 4.1 (The Sieve Weight Quotient). *To evaluate the structural efficiency of a given weight function, we define the sieve weight quotient $\mu(n, W)$ as the ratio of the penalized weight sum to the total weight sum:*

$$\mu(n, W) = \frac{S_2(n, W)}{S_1(n, W)} \quad (2)$$

By Theorem 3.2, if we can construct a weight function such that $\mu(n, W) < 1$, the sieve sufficiency condition (Definition 3.3) is satisfied.

Having isolated the destructive interference at $h = 3q/p_i$, a natural first attempt is to construct weight functions composed entirely of independent, one-dimensional exponential sums. We can assess the limitations of this approach by examining the isolated spectral capacity of the sieve.

Definition 4.2 (Principal Density and Character Suppression Capacity [43 & 44]). *For a one-dimensional character weight system evaluated over \mathcal{P}_n , the principal expected density is governed by the harmonic sum over the prime window:*

$$\text{Principal Term}(n) = \sum_{p_i \in \mathcal{P}_n} \frac{2}{p_i}$$

The corresponding maximum theoretical suppression available from the targeted frequencies is defined as the suppression capacity:

$$\text{Suppression Capacity}(n) = \sum_{p_i \in \mathcal{P}_n} \frac{2}{p_i} \cos\left(\frac{\pi}{p_i}\right)$$

If the suppression capacity could exceed or match the principal density term, a one-dimensional sieve could theoretically suppress the penalized sum S_2 below S_1 . However, the geometry of the character sum strictly forbids this.

Lemma 4.3 (The Spectral Deficit [44]). *For any valid prime window \mathcal{P}_n , the suppression capacity of any purely one-dimensional character weight system is strictly less than the sieve's principal density term:*

$$\sum_{p_i \in \mathcal{P}_n} \frac{2}{p_i} \cos\left(\frac{\pi}{p_i}\right) < \sum_{p_i \in \mathcal{P}_n} \frac{2}{p_i} \quad (3)$$

Proof. For any prime $p_i \geq 5$, the phase angle satisfies $0 < \pi/p_i < \pi/2$, which strictly bounds the cosine function: $\cos(\pi/p_i) < 1$. Multiplying by the strictly positive amplitude $2/p_i$ yields

$$\frac{2}{p_i} \cos\left(\frac{\pi}{p_i}\right) < \frac{2}{p_i}$$

Summing these strict, term-by-term inequalities across the finite non-empty set \mathcal{P}_n preserves the strict inequality:

$$\sum_{p_i \in \mathcal{P}_n} \frac{2}{p_i} \cos\left(\frac{\pi}{p_i}\right) < \sum_{p_i \in \mathcal{P}_n} \frac{2}{p_i}$$

which completes the proof. \square

While the inequality itself is algebraically elementary, its significance lies in the asymptotic behavior of the divergence.

Proposition 4.4 (One-Dimensional Parity Barrier). *For any sufficiently large n , utilizing purely independent, one-dimensional exponential sum weights forces the sieve weight quotient to $\mu(n, W) \geq 1$: it is structurally impossible for this geometrically symmetric additive sieve to resolve the twin prime conjecture using solely one-dimensional variance bounds.*

Proof. By Mertens' theorems [14], the principal density term $\sum \frac{2}{p_i}$ diverges logarithmically as $n \rightarrow \infty$. Conversely, the magnitude of the spectral deficit (3), measured by the difference between the principal term and the suppression capacity, can be bounded using the Taylor expansion $\cos(x) \geq 1 - x^2/2$:

$$\text{Deficit} = \sum_{p_i \in \mathcal{P}_n} \frac{2}{p_i} \left(1 - \cos \frac{\pi}{p_i}\right) \leq \sum_{p_i \in \mathcal{P}_n} \frac{\pi^2}{p_i^3}$$

Because the sum $\sum p^{-3}$ converges over the primes, the absolute suppression capacity of the one-dimensional character evaluations grows slower than the logarithmically diverging principal density. However, establishing the strict condition $\mu(n, W) \geq 1$ requires examining the quotient S_2/S_1 . The suppression sum relies on the specific Fourier coefficients $\hat{W}(3q/p_i)$. For a purely one-dimensional weight function $W(x) \geq 0$, these coefficients are constrained by the total weight $q\hat{W}(0)$. Since $W(x)$ cannot simultaneously saturate the theoretical suppression capacities at $\hat{W}(3q/p_i)$ for all p_i without taking negative values, the actual achievable suppression is strictly bounded away from the theoretical maximum. Consequently, the logarithmically growing gap between the principal density and the available one-dimensional suppression guarantees that for all sufficiently large n , the ratio $\mu(n, W)$ cannot be suppressed below 1. \square

This convergence-divergence disparity forces the quotient $\mu(n, W) \geq 1$ for any sufficiently large n when utilizing one-dimensional independent evaluations. This spectral deficit is the analytic manifestation of Selberg's parity barrier within our additive framework. (A finite truncation of purely independent character sums will inevitably force the weight function to take negative values at dense overlapping forbidden regions before it can reduce the overall weight quotient below 1.)

This obstruction shows that to circumvent the density deficit and overcome the parity barrier, the sieve weight construction must incorporate multidimensional exponential sums and higher-order correlations.

5. VARIATIONAL OPTIMIZATION OF THE SIEVE WEIGHT QUOTIENT

The analytic barrier established in Section 4 shows that utilizing purely independent one-dimensional exponential sums fundamentally fails to bypass the structural density deficit. To constructively overcome this, we must transition to a correlated, multidimensional weight space. We build this space using higher-order correlations of the fundamental character variables.

Definition 5.1 (Multidimensional Basis and Polynomial [↵ & ↵]). *For any subset of prime indices $S \subseteq \{1, \dots, w(n)\}$, we define the multidimensional basis function as the product of the targeted trigonometric evaluations:*

$$B_S(x) = \prod_{i \in S} \cos\left(\frac{6\pi x}{p_i}\right)$$

Let λ be a real-valued coefficient vector indexed by the power set of $\{1, \dots, w(n)\}$. We construct the multidimensional polynomial $P(x)$ as a linear combination of these basis functions:

$$P(x) = \sum_S \lambda_S B_S(x)$$

The resulting $2^{w(n)}$ -dimensional basis is closed under pointwise multiplication, which allows us to ensure the non-negativity of the sieve weight function.

Definition 5.2 (The Multidimensional Weight Function [43]). *We define the multidimensional weight function $W_\lambda(x)$ by squaring the polynomial $P(x)$ and restricting its support to the evaluation interval \mathcal{A}_n :*

$$W_\lambda(x) = \begin{cases} (P(x))^2 & \text{if } x \in \mathcal{A}_n \\ 0 & \text{otherwise} \end{cases}$$

By casting the weight space as the square of a polynomial, we automatically satisfy the condition $W_\lambda(x) \geq 0$ for all x , avoiding the need to impose infinitely many linear inequality constraints on λ . This formulation translates our sieve sums into classical quadratic forms.

Lemma 5.3 (Reduction to Quadratic Forms [43 & 44]). *The total weight sum $S_1(n, W_\lambda)$ and the penalized weight sum $S_2(n, W_\lambda)$ map directly to continuous quadratic forms $Q_1(\lambda)$ and $Q_2(\lambda)$. Specifically, the total weight sum evaluates to the squared norm of the polynomial over the interval:*

$$S_1(n, W_\lambda) = Q_1(\lambda) = \sum_{x \in \mathcal{A}_n} (P(x))^2$$

Proof. By definition, $S_1 = \sum_{x \in \mathcal{A}_n} W_\lambda(x)$. Substituting the construction $W_\lambda(x) = (P(x))^2$, we expand the square and exchange the order of summation:

$$\begin{aligned} S_1(n, W_\lambda) &= \sum_{x \in \mathcal{A}_n} \left(\sum_S \lambda_S B_S(x) \right)^2 \\ &= \sum_{x \in \mathcal{A}_n} \sum_{S, T} \lambda_S B_S(x) B_T(x) \lambda_T \\ &= \sum_{S, T} \lambda_S \left(\sum_{x \in \mathcal{A}_n} B_S(x) B_T(x) \right) \lambda_T \end{aligned}$$

Defining the matrix elements $M_1(S, T) = \sum_{x \in \mathcal{A}_n} B_S(x) B_T(x)$ yields the quadratic form $Q_1(\lambda) = \sum_{S, T} \lambda_S M_1(S, T) \lambda_T$.

The derivation for $S_2(n, W_\lambda)$ proceeds identically by incorporating the linear penalty function $c(x)$ into the inner sum:

$$\begin{aligned} S_2(n, W_\lambda) &= \sum_{x \in \mathcal{A}_n} c(x) \left(\sum_S \lambda_S B_S(x) \right)^2 \\ &= \sum_{S, T} \lambda_S \left(\sum_{x \in \mathcal{A}_n} c(x) B_S(x) B_T(x) \right) \lambda_T \end{aligned}$$

yielding the quadratic form $Q_2(\lambda)$. \square

With both sums expressed as quadratic forms, the sieve weight quotient (2) transitions to a Rayleigh-Ritz quotient $R(\lambda) = Q_2(\lambda)/Q_1(\lambda)$ [43]. Because this quotient is inherently scale-invariant, we can apply a standard compactness argument.

Lemma 5.4 (Compactness of the Orthogonal Unit Sphere [43]). *Let $\ker(Q_1)$ denote the kernel of the quadratic form Q_1 . Because the quotient is invariant under both uniform scaling and the additive translation of vectors from $\ker(Q_1)$, we restrict the parameter space to the unit sphere within the orthogonal complement, denoted as \mathcal{S}^\perp . This bounded subspace \mathcal{S}^\perp is compact.*

Proof. Within the finite-dimensional Euclidean space of the coefficients, compactness is equivalent to the subspace being both closed and bounded (Heine-Borel theorem). The orthogonal complement $\ker(Q_1)^\perp$ forms a closed linear subspace. The unit sphere condition, enforced by the standard dot product $\lambda \cdot \lambda = 1$, defines a set that is closed and bounded. Because \mathcal{S}^\perp is the intersection of these sets, it is compact. \square

Having established Lemma 5.4, we can now analyze the continuous image of this subspace under the quotient mapping. This allows us to define the absolute lower bound of the sieve weight quotient (2).

Definition 5.5 (The Minimum Attainable Ratio). *We define the set of attainable ratios as the continuous image of the compact space \mathcal{S}^\perp under the mapping $R(\lambda)$. The minimum sieve quotient $\mu_{\min}(n)$ is defined as the absolute infimum of this attainable set.*

By using the topological compactness derived above, we can assert the existence of an optimal coefficient vector.

Theorem 5.6 (Existence of the Exact Minimizer [43]). *The minimum sieve quotient $\mu_{\min}(n)$ is attained by an optimal coefficient vector $\lambda_{\text{opt}} \in \mathcal{S}^\perp$, such that $Q_1(\lambda_{\text{opt}}) > 0$ and $R(\lambda_{\text{opt}}) = \mu_{\min}(n)$.*

Proof. The quotient $R(\lambda)$ is continuous wherever $Q_1(\lambda) > 0$. By confining our domain to the orthogonal complement \mathcal{S}^\perp , we exclude the kernel of Q_1 , ensuring continuity across the entire restricted domain. Because continuous functions map compact spaces to compact spaces, the resulting set of attainable ratios is itself compact. By the extreme value theorem, a compact subset of the real numbers contains its infimum, guaranteeing the existence of the minimizer λ_{opt} . \square

With the existence of the minimiser established, we obtain the main result of this section. We can now link the continuous bound $\mu_{\min}(n)$ back to the discrete existence of twin prime indices within our evaluation interval.

Theorem 5.7 (Optimal Sufficiency and Sieve Guarantee [43 & 43]). *The optimal multidimensional weight function $W_{\text{opt}}(x)$ constructed from λ_{opt} satisfies the sufficiency condition ($S_2 < S_1$) if and only if $\mu_{\min}(n) < 1$. If this analytic bound holds, the framework guarantees the existence of an index for a twin prime pair within \mathcal{A}_n .*

Proof. By the definition of the minimum ratio, evaluating the moments with the optimal weight function yields exactly $S_2(n, W_{\text{opt}})/S_1(n, W_{\text{opt}}) = R(\lambda_{\text{opt}}) = \mu_{\min}(n)$. Since S_1 is positive for the optimal non-zero weight, it follows algebraically that $S_2 < S_1$ holds if and only if the ratio $\mu_{\min}(n) < 1$. If this inequality is satisfied, the weight function W_{opt} fulfills the sufficiency condition of Definition 3.3. By directly applying Theorem 2.4, this sufficiency implies the existence of at least one index $x \in \mathcal{A}_n$ that generates a twin prime pair. \square

By combining the per-interval guarantee with a standard infinitary argument, we obtain the logical capstone of the framework.

Theorem 5.8 (Conditional Twin Prime Theorem [44]). *If there exist infinitely many integers $n \geq 1$ for which $\mu_{\min}(n) < 1$, then there are infinitely many twin primes.*

Proof. Suppose the set $\{n \in \mathbb{N} \mid \mu_{\min}(n) < 1\}$ is infinite. For any bound B , choose n in this set with $n > B$. By Theorem 3.4, there exists an index $x \in \mathcal{A}_n$ such that $6x - 1$ and $6x + 1$ are both prime. Since $x \geq 6n^2 - 2n \geq 4n > 4B$, the smaller prime satisfies $6x - 1 > B$. As B was arbitrary, the set of primes p such that both p and $p + 2$ are prime is unbounded, hence infinite. \square

This theorem completes the reduction. The parity obstruction has been translated into a concrete analytic condition: the twin prime conjecture follows if one can establish that the spectral density of the higher-order correlations forces $\mu_{\min}(n) < 1$ for infinitely many n .

6. DISCUSSION: HIGHER-ORDER CORRELATIONS AND FUTURE DIRECTIONS

While the reduction of the twin prime conjecture to the spectral bound $\mu_{\min}(n) < 1$ provides a rigid finite-dimensional framework, achieving this bound analytically remains obstructed by the spectral deficit. However, the multidimensional parameter space \mathcal{S}^\perp theoretically bypasses this limitation. We can map the analytic mechanism of this bypass by examining the full inclusion-exclusion convolution in the frequency domain.

6.1. The perfect sieve weight and character convolution. To understand how higher-order correlations circumvent the parity barrier, consider the "perfect" spatial weight function $\omega(x)$, which equals 1 when $c(x) = 0$ and 0 elsewhere. It is given by the pointwise product of the allowed residue indicators:

$$\omega(x) = \prod_{i=1}^{w(n)} (1 - g_i(x))$$

Transforming this identity into the frequency domain over $\mathbb{Z}/q\mathbb{Z}$ yields a spectral characterization. Let $\hat{g}_i(h)$ be the Fourier transform of the individual indicator functions. By the convolution theorem, the sequence of allowed residues transforms as $q\delta_{h,0} - \hat{g}_i(h)$. Consequently, the perfect multidimensional weight function in the frequency domain resolves into the full character convolution:

$$\Omega_h = (q\delta_{h,0} - \hat{g}_1(h)) \circledast (q\delta_{h,0} - \hat{g}_2(h)) \circledast \cdots \circledast (q\delta_{h,0} - \hat{g}_{w(n)}(h))$$

where \circledast denotes the circular discrete convolution over $\mathbb{Z}/q\mathbb{Z}$.

When this convolution is expanded, it naturally generates higher-order cross-correlations (e.g., $\hat{g}_i(h_1) \circledast \hat{g}_j(h_2)$). These multidimensional exponential sums appear to be indispensable. When transitioning to a higher-dimensional polynomial basis space, the resulting interference cross-correlations intuitively act as corrective phases. We expect them to constructively interfere at the intersections of forbidden congruences, forcing the weight function to vanish. This multidimensional scaling is thus intended to simultaneously preserve non-negativity and protect the principal density term from over-subtraction, an effect which is empirically supported by numerical optimization of the quotient for small values of n .

6.2. Numerical explorations of higher-order correlations. To empirically validate the structural necessity of these multidimensional corrective phases, we implemented the continuous variational framework and computationally minimized the sieve quotient $\mu_{\min}(n)$ via the generalized eigenvalue problem derived from the discrete penalty function. The Python scripts developed for these spectral optimizations are publicly available in the accompanying GitHub repository.

Due to the exponential growth of the basis set of multidimensional sieve weights, a direct calculation incorporating all prime factors and unrestricted correlation depths rapidly exceeds standard computational resources. Therefore, our initial investigations restricted the optimization to two-dimensional correlations (evaluating interference strictly between pairs of primes). Under this restriction, we obtained admissible quotients $\mu_{\min}(n) < 1$ for all bounded ranges up to $n = 260$. Figure 5 illustrates the trajectory of the optimized 2D quotient. While the 2D correlations initially provide sufficient destructive interference to suppress the quotient below 1, the distinct upward trend suggests that the spectral deficit induced by the one-dimensional Mertens terms ultimately overcomes strictly pairwise interactions as n grows.

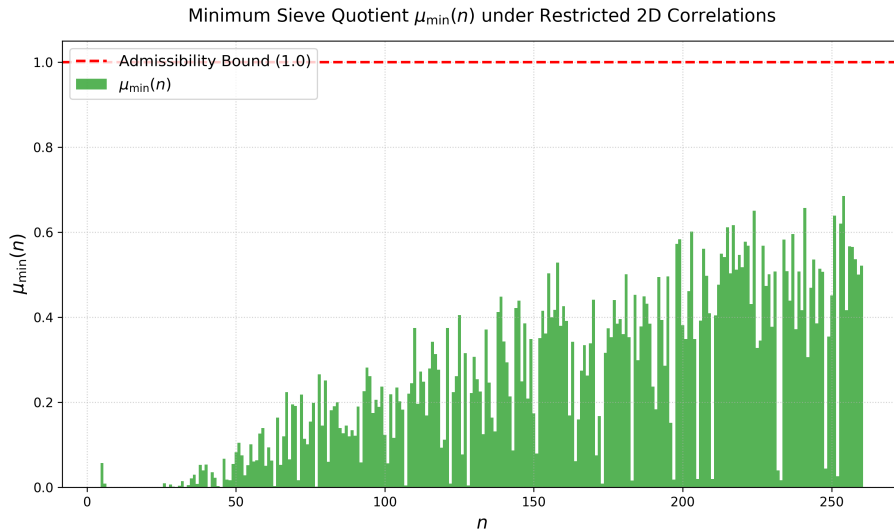


FIGURE 5. The minimum sieve quotient $\mu_{\min}(n)$ under restricted 2D correlations. The upward trend indicates that pairwise interference struggles to maintain the admissibility bound $\mu < 1$ for larger n .

To establish the essential role of higher-order interactions in permanently breaching the parity barrier, we performed a targeted experiment at $n = 199$. In the purely 2D model utilizing the complete prime basis for $n = 199$, the optimization yields a quotient of $\mu_{\min}(199) \approx 0.583$.

We subsequently relaxed the correlation depth to 3D, while restricting the basis to a highly active subset of 20 prominent primes (to keep the matrix dimension computationally feasible at $D = 1351$). The inclusion of 3D correlations sharply reduced the minimum quotient to $\mu_{\min}(199) \approx 0.081$. This sharp reduction corroborates the

core analytic thesis of this paper: while 1D and 2D components struggle against the logarithmic divergence of the local penalty, higher-dimensional cross-correlations provide sufficiently large negative feedback to permanently satisfy the condition $\mu_{\min}(n) < 1$.

6.3. Towards bounding the minimum ratio. The exact resolution of this convolution computationally requires navigating the massive $2^{w(n)}$ -dimensional polynomial space, making direct computation infeasible for large n . However, establishing the Sieve Sufficiency Guarantee unconditionally only requires demonstrating that $\mu_{\min}(n) < 1$. By the continuous variational principle, we are not required to locate the exact optimal minimizer λ_{opt} . Any analytically constructed test vector λ_{test} corresponding to a weight matrix that pushes the sieve weight quotient strictly below 1 is sufficient.

Advancing the resolution of the twin prime conjecture within this additive sieve framework therefore requires the following analytical steps:

- (1) Expand the theoretical multidimensional convolution Ω_h and extract the lowest-order, highest-magnitude cross-correlation coefficients.
- (2) Map these theoretical phase corrections onto the finite-dimensional polynomial basis $B_S(x) = \prod_{i \in S} \cos(6\pi x/p_i)$.
- (3) Construct a continuous bounded weight vector λ_{test} over the target interval \mathcal{A}_n .
- (4) Analytically bound the quadratic limits $Q_2(\lambda_{\text{test}})/Q_1(\lambda_{\text{test}})$, proving that the scaled cross-correlations possess the capacity to overcome the diverging spectral deficit.

While executing this final step demands careful estimation of the resulting combinatorial sums, the framework nevertheless provides a well-defined approach. Importantly, this continuous analytic framework maps a combinatorial obstacle into an explicit, seemingly resolvable optimization problem.

7. CONCLUSION

This paper has presented a rigorous framework that equates the existence of twin primes to the bounded optimization of a specific residue structure. By tracking localized sieve penalties over symmetric bounded intervals and translating them into the frequency domain via the discrete Fourier transform, we mapped discrete prime alignments to continuous character sum spectra.

This synthesis of classical techniques—ranging from Diophantine forms and Turán variances to Plancherel’s theorem and topological compactness—provides a spectral lens on the twin prime problem. Crucially, it allows us to isolate the exact analytic manifestation of the parity barrier. We proved that one-dimensional exponential sum weights suffer from a structural deficit against the logarithmic divergence of Mertens’ sums, forcing the sieve quotient $\mu(n, W) \geq 1$ and formally necessitating the use of multidimensional polynomial bases.

Ultimately, our reduction to $\mu_{\min}(n) < 1$ replaces the classical combinatorial parity obstruction with a concrete spectral bound. The problem is thus reframed entirely into estimating the spectral density of heavily structured higher-order cross-correlations. For researchers exploring this domain, this explicit variational framework provides a novel avenue for approaching the parity barrier. The analytical

bounding of these multidimensional correlation limits remains a challenging, yet now explicitly defined, open problem.

AI USAGE DISCLOSURE

The author acknowledges the use of generative AI (Gemini 3.1 Pro [9]) for LaTeX technical synthesis and structural formatting, as well as for structuring the Lean 4 formalization by designing precise prompt specifications for implementation by the Aristotle [1] automated theorem prover. The core mathematical logic, proofs, and final validation were performed by the author.

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